

Quantum computing for seismic problems

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Outline

1. Background and motivation

2. Quantum computing

- a) Quantum bit (qubit)
- b) Quantum parallelism
- c) Physical realization of qubit
- d) Mathematics of qubit

3. Applications

- a) Quantum algorithms
- b) Seismic algorithms

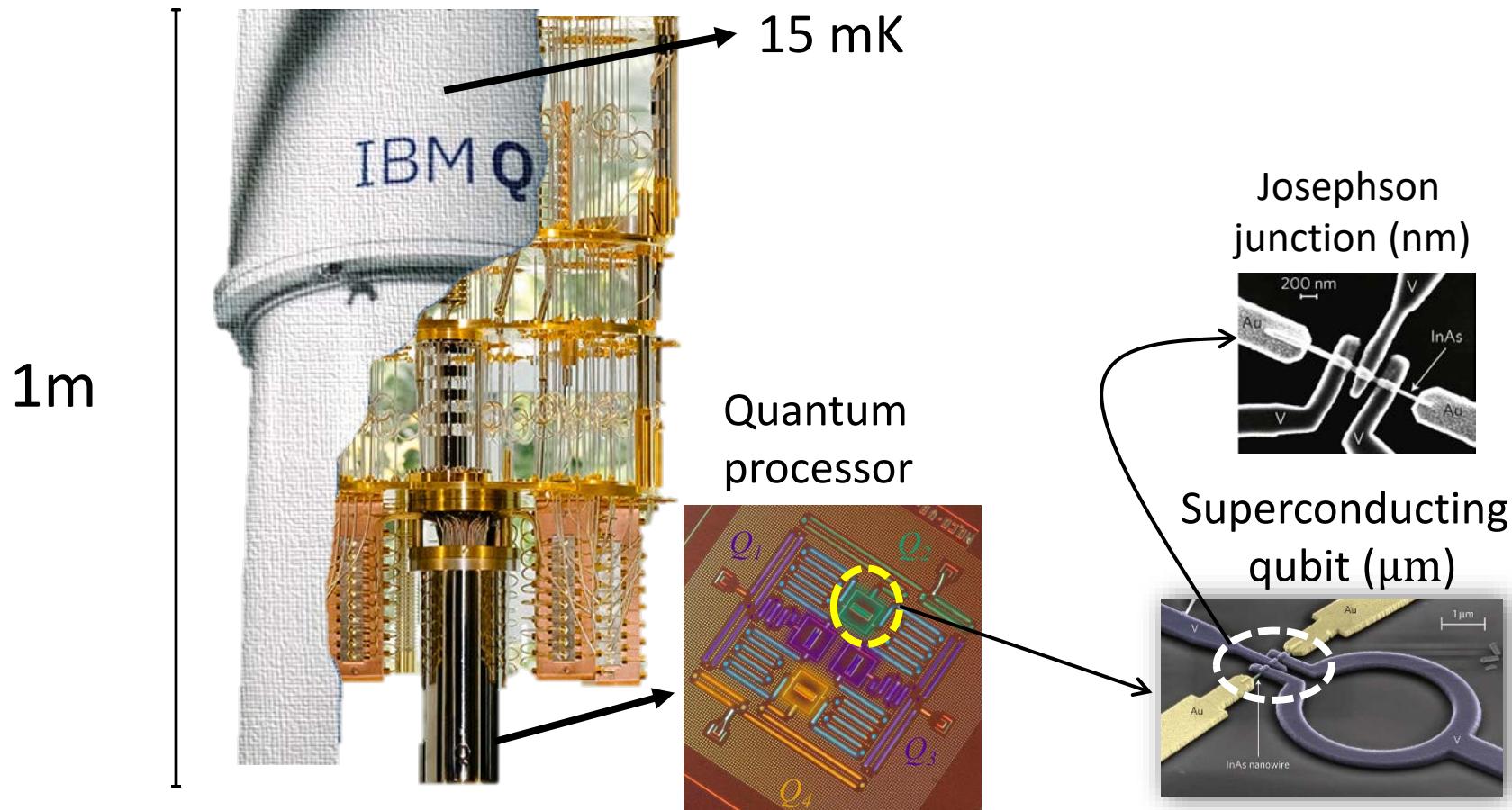
4. Summary and future works

Motivation

- Quantum algorithms for seismic wave modeling
- Develop the software package for quantum computing with applications in seismic problems
- Design a quantum simulator for modeling and inversion

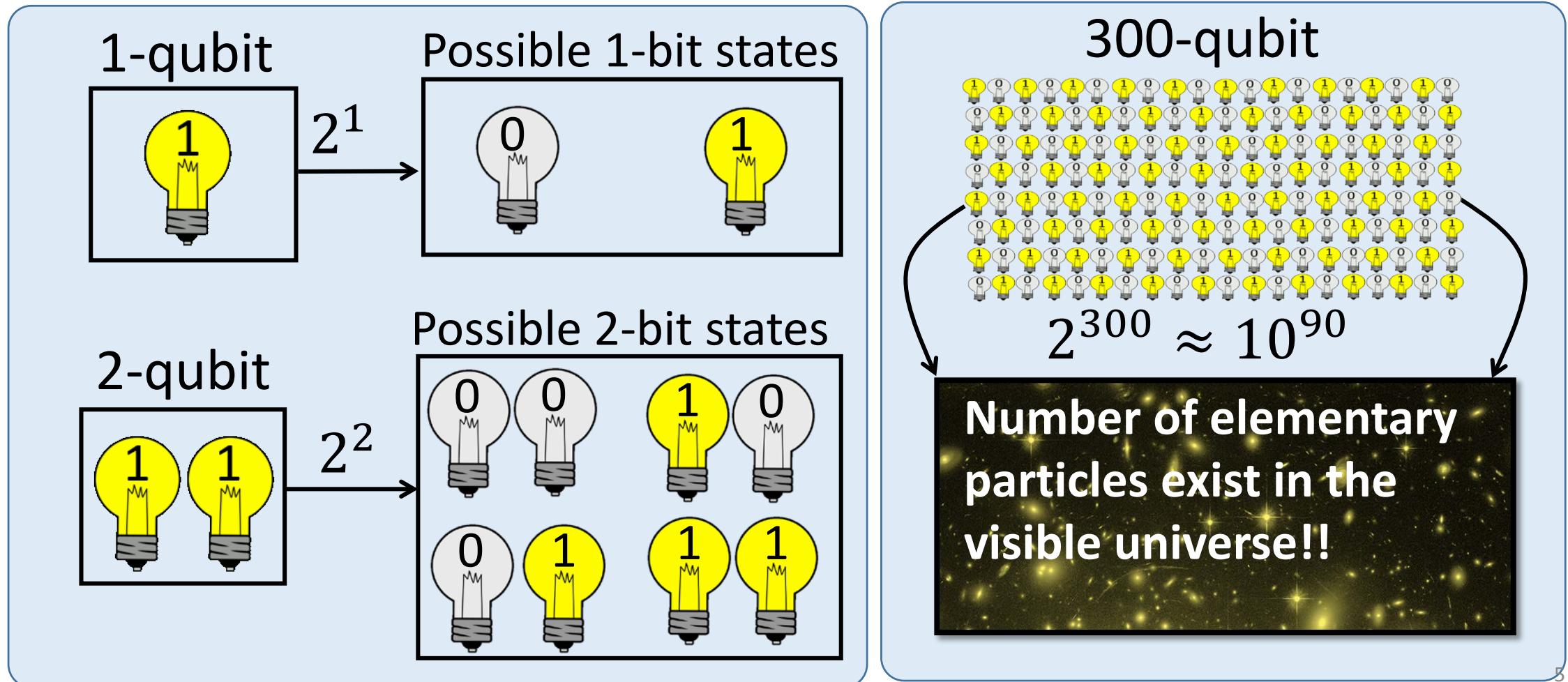
Background

IBMQ (2017) 50 qubit quantum computer



Quantum computing

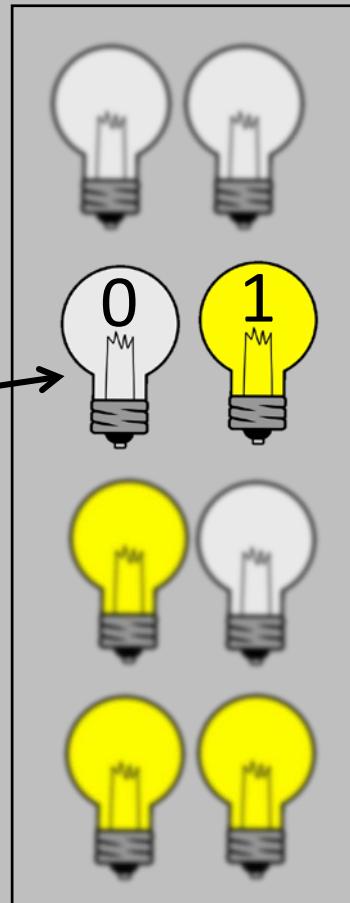
qubit: the unit of quantum information



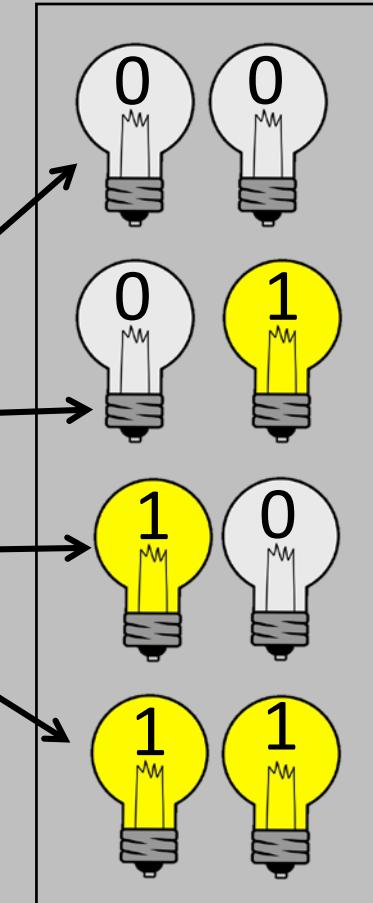
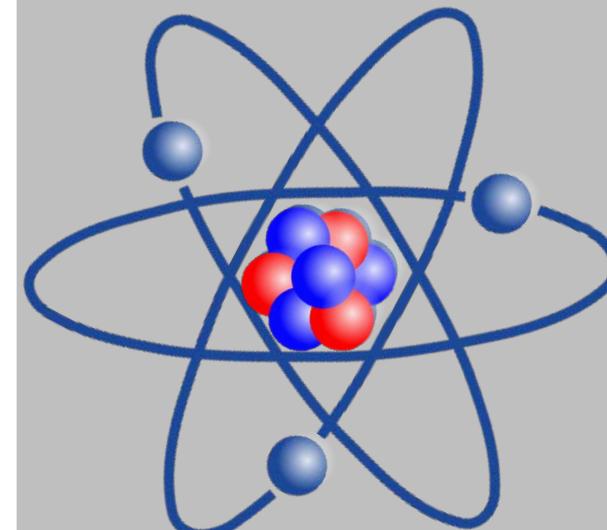
Quantum computing

Quantum parallelism: computational speedup

2-bit
Classical computer



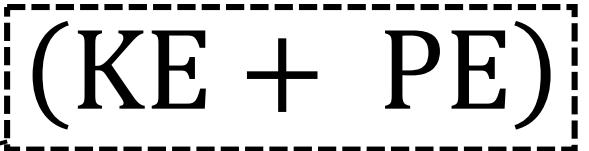
2-qubit
Quantum computer



Quantum computing

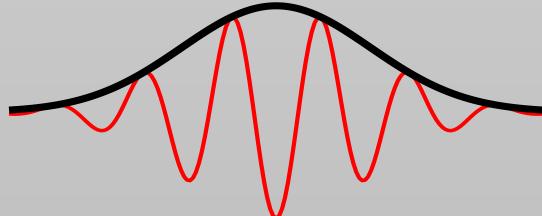
Schrodinger equation: wave equation for atoms

$$[(KE + PE)] \Psi = E\Psi$$

H(Hamiltonian) 

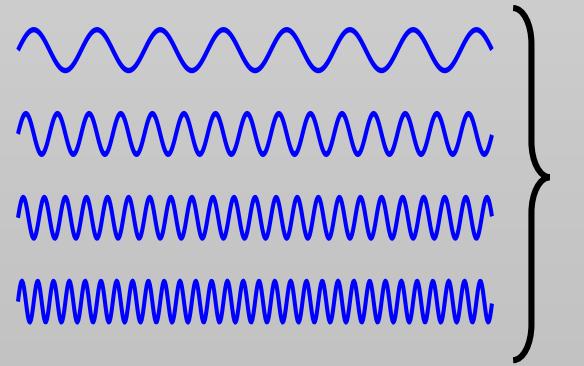
Algorithm output

$|\Psi(x, t)|^2$ Probability density
for finding the particle at
position X and time t



Quantum parallelism

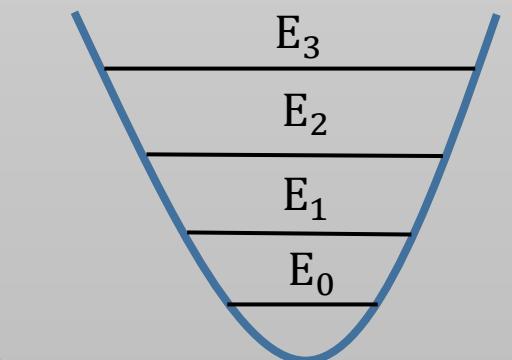
Superposition principle



Hardware (Q proc.)

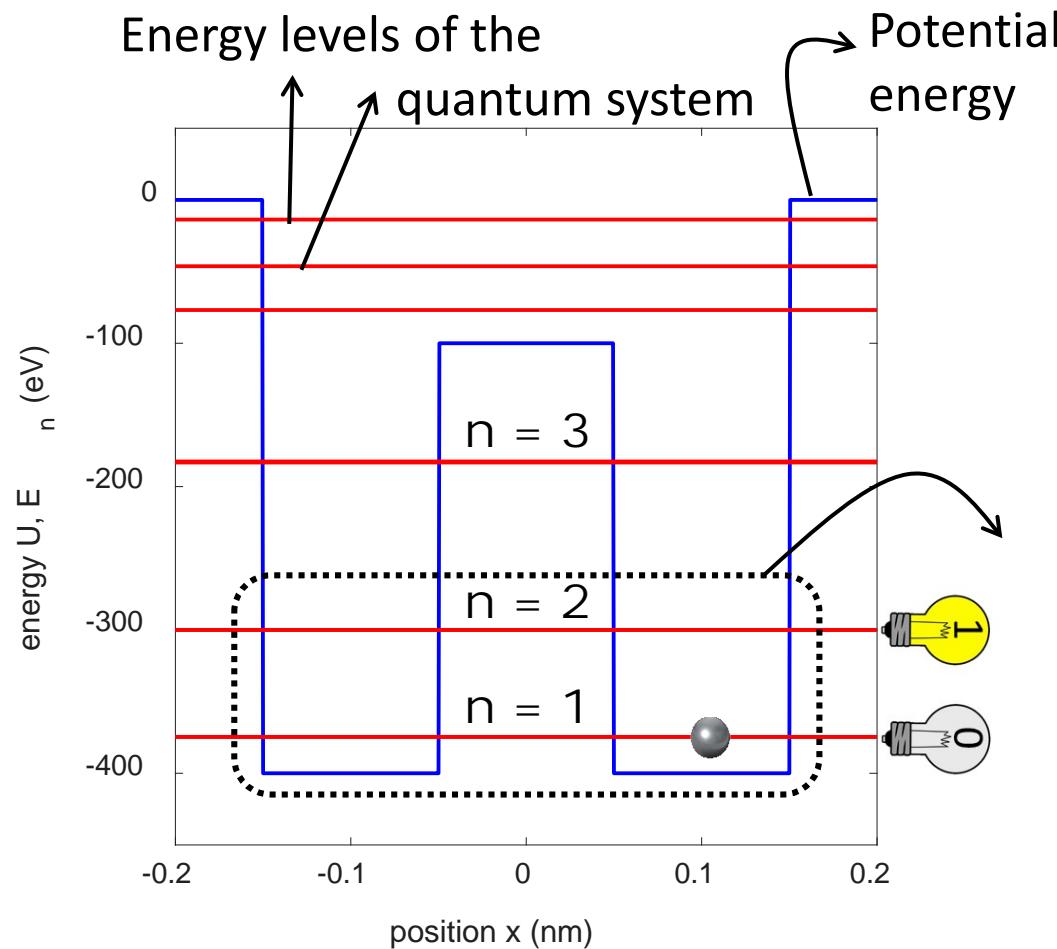
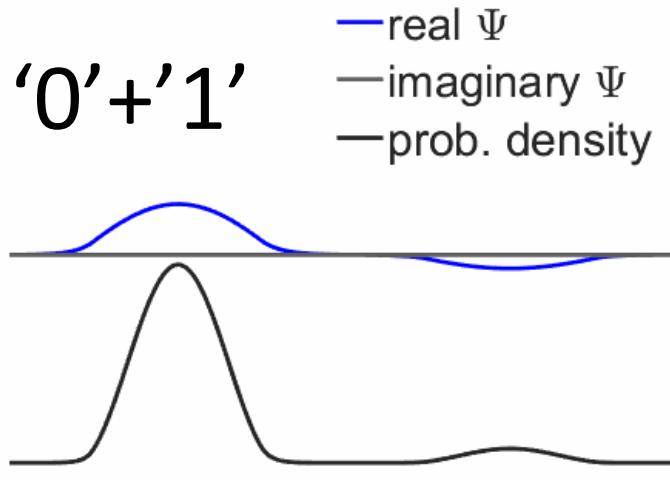
Energy is discretized

$$H\Psi = E\Psi$$

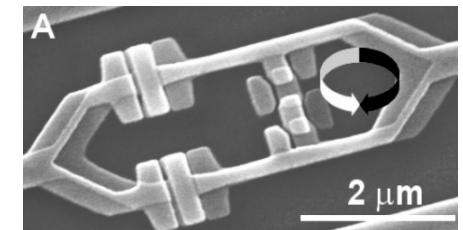


Quantum computing

Physical simulation of qubit



Superconducting phase qubit



Qubit is the superposition of the lowest energy levels of the system near absolute zero temperature

Quantum computing

Mathematics of qubit

Ket notation

$$0 \rightarrow |0\rangle \equiv \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

$$1 \rightarrow |1\rangle \equiv \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$

$$1\text{-qubit } |\Psi\rangle = \alpha|0\rangle + \beta|1\rangle$$

$$|\alpha|^2$$

Probability that
qubit to be in $|0\rangle$

$$|\beta|^2$$

Probability that
qubit to be in $|1\rangle$

Quantum computing

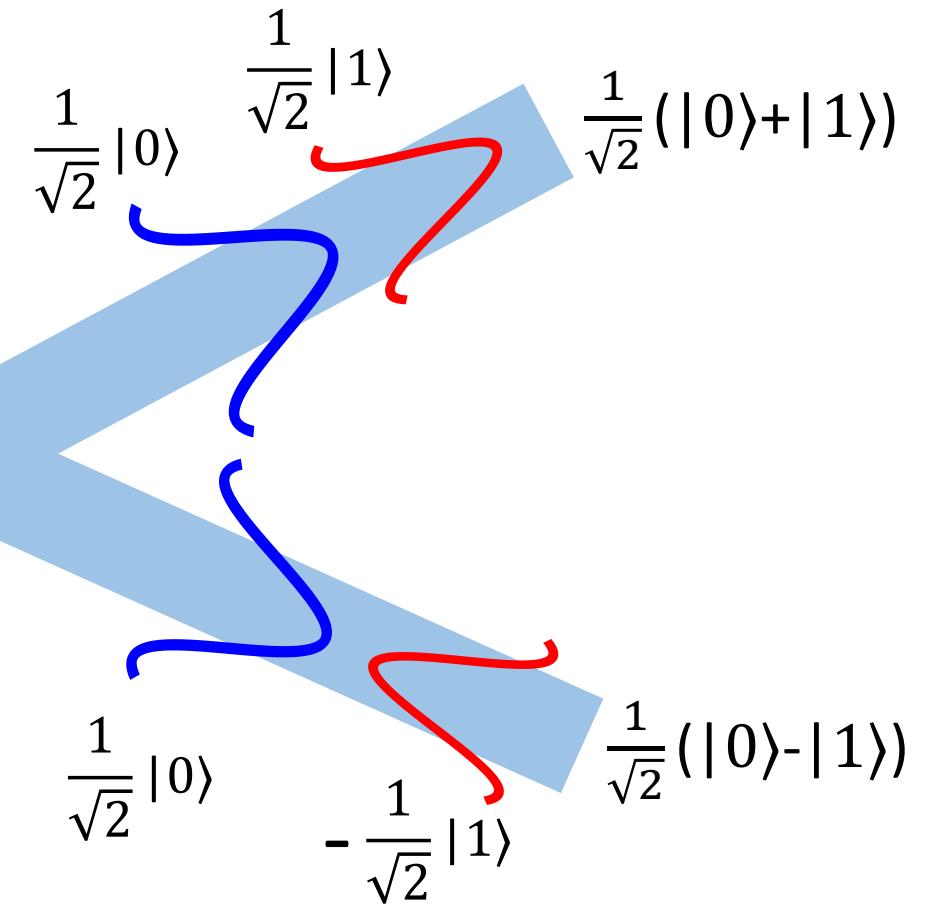
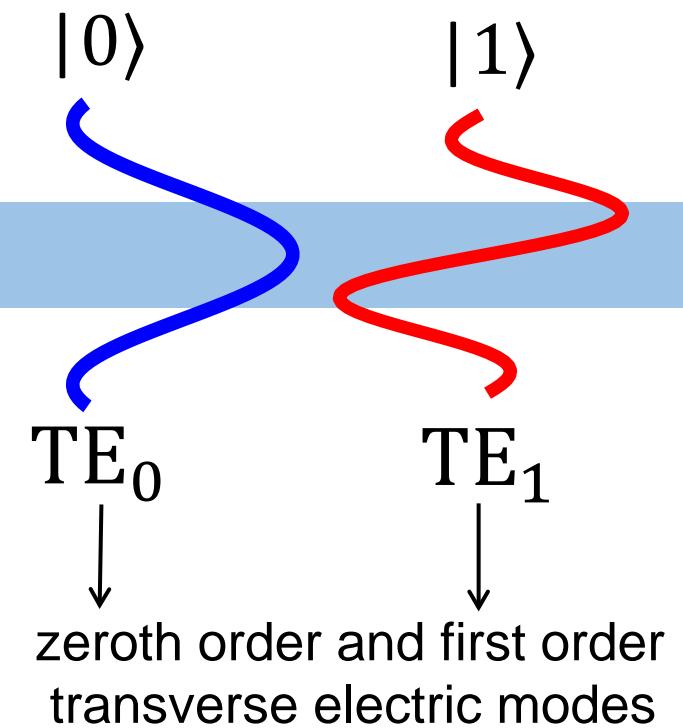
Quantum gate

Hadamard gate

$$H|0\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$$

$$H|1\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

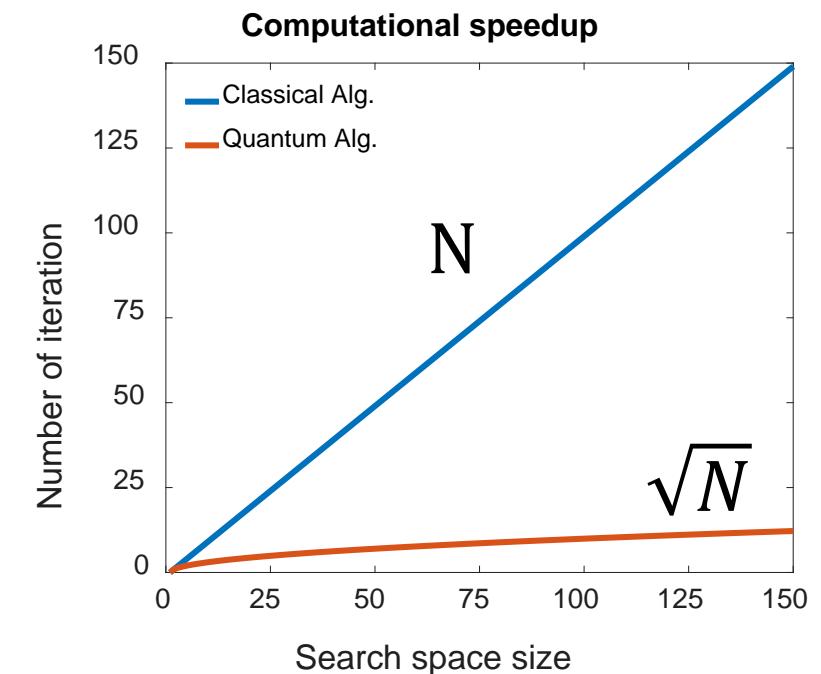
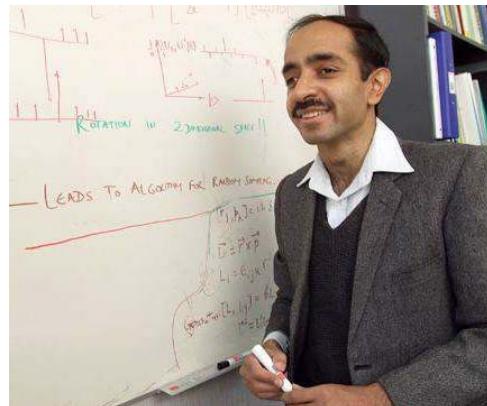
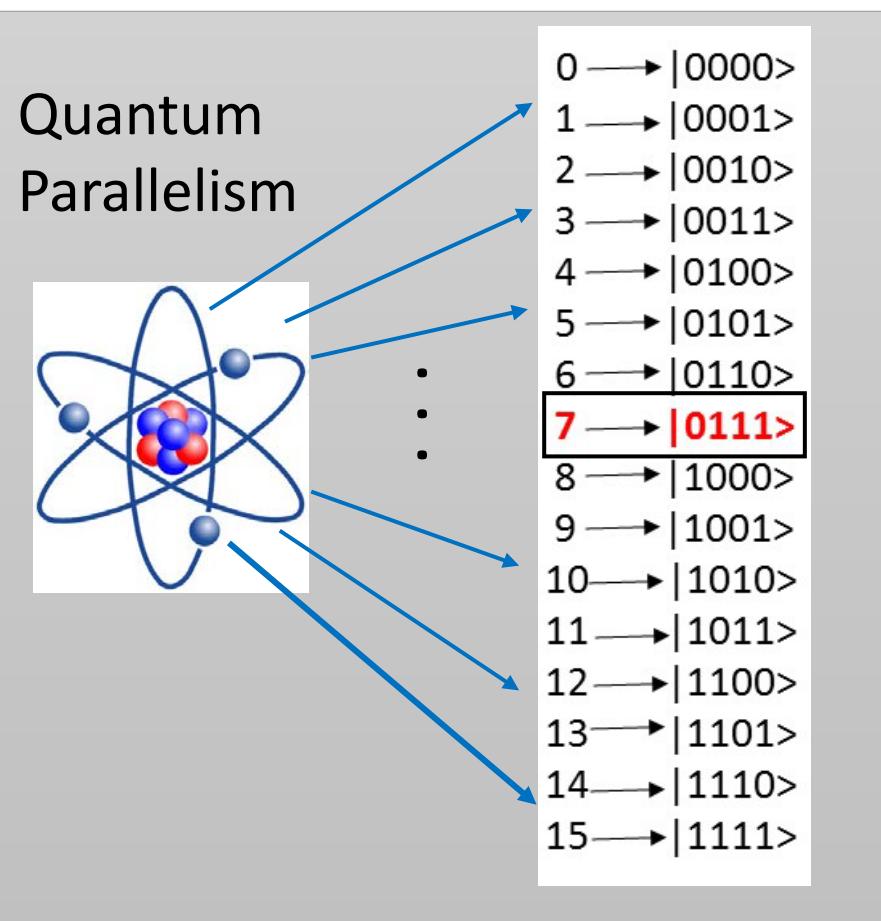
Y-junction
beam splitter



Applications

Quantum search algorithm

Quantum mechanics helps in searching for a needle in a haystack- Grover-PRL, 1997

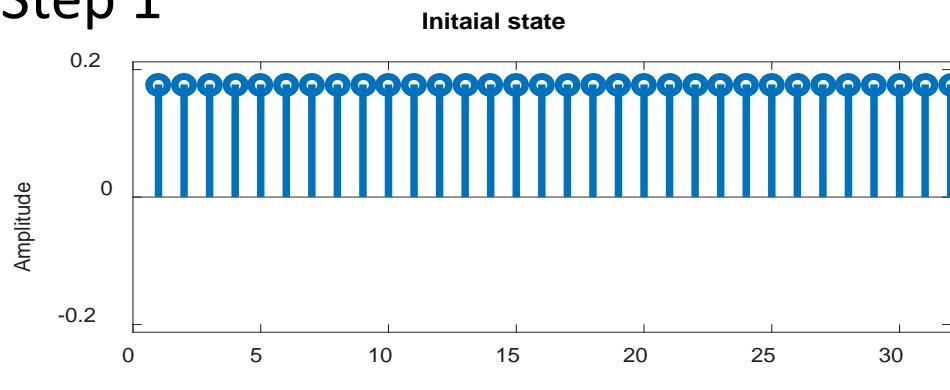


Applications

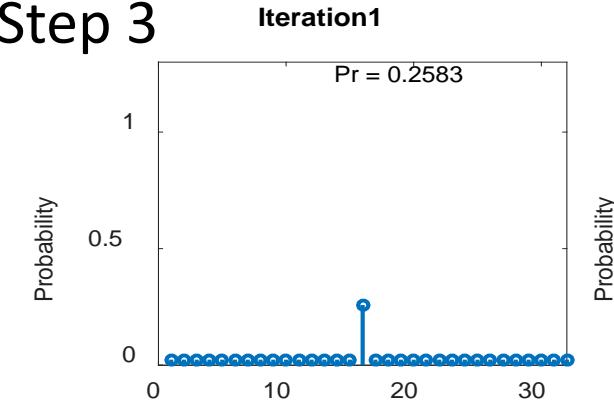
Quantum search algorithm

Numerical modeling $N=32=2^5$, simulation with 5 qubits

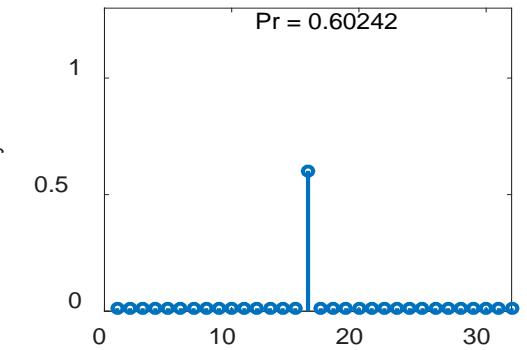
Step 1



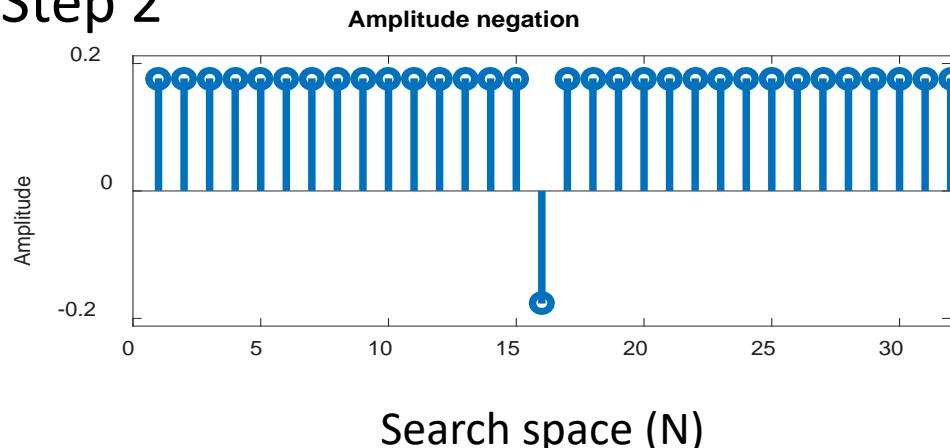
Step 3



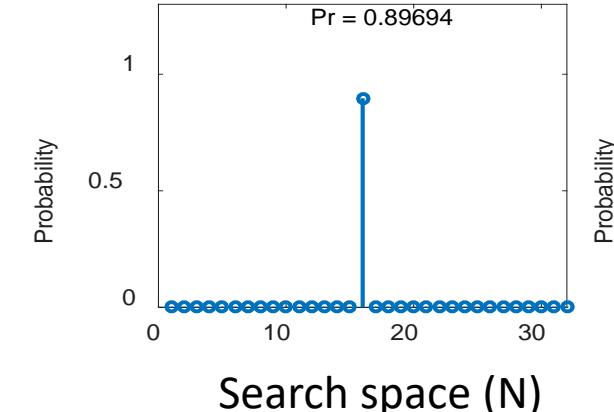
Iteration2



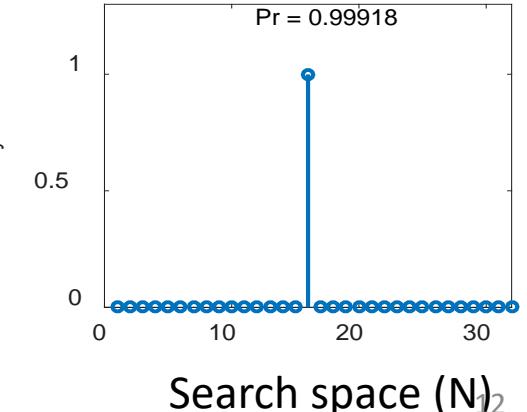
Step 2



Iteration3



Iteration4



Applications

Quantum Fourier Transform

DFT

DFT: $(x_0, \dots, x_{N-1}) \mapsto (y_0, \dots, y_{N-1})$

$$y_k = \frac{1}{\sqrt{N}} \sum_{j=0}^{N-1} x_j e^{i2\pi jk/N}$$

DFT on $N = 2^n$ elements requires $O(n2^n)$ gates

QFT

$$\text{QFT: } \sum_{j=1}^{N-1} x_j |j\rangle \mapsto \sum_{k=1}^{N-1} y_k |k\rangle$$

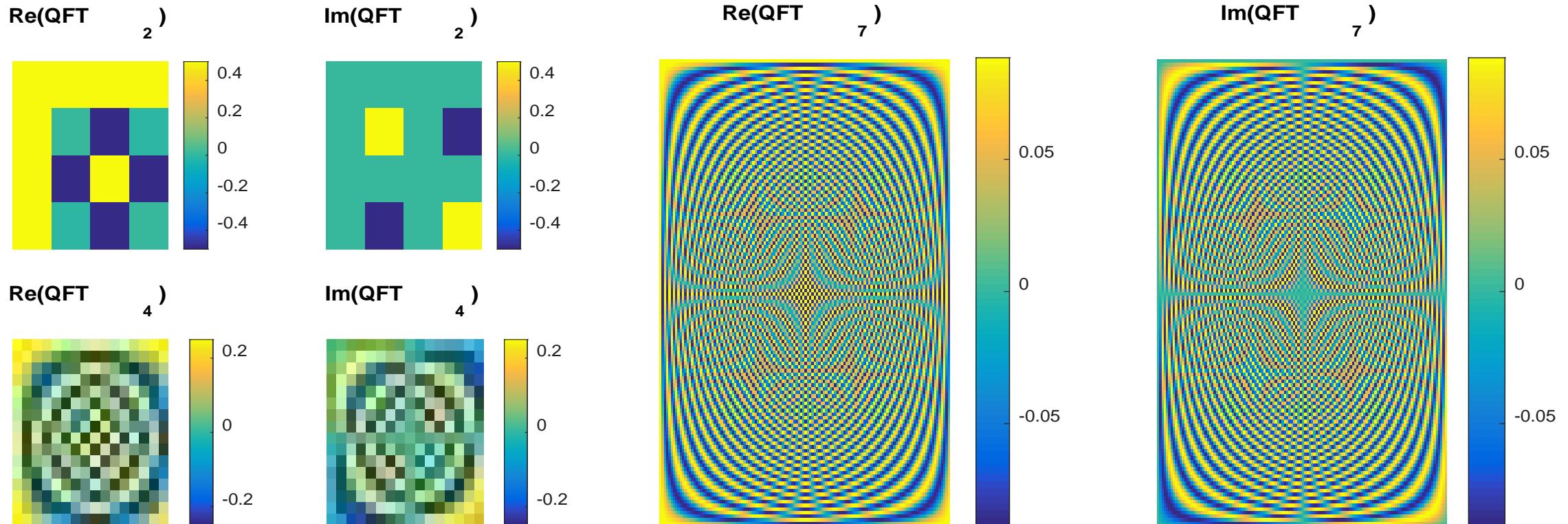
$$|j\rangle \mapsto \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} e^{i2\pi jk/N} |k\rangle$$

QFT on $N = 2^n$ elements requires only $O(n^2)$ gates (**Exponential speed up!**)

Applications

Quantum Fourier Transform

Numerical modeling



Applications

Quantum algorithm for linear systems of equations

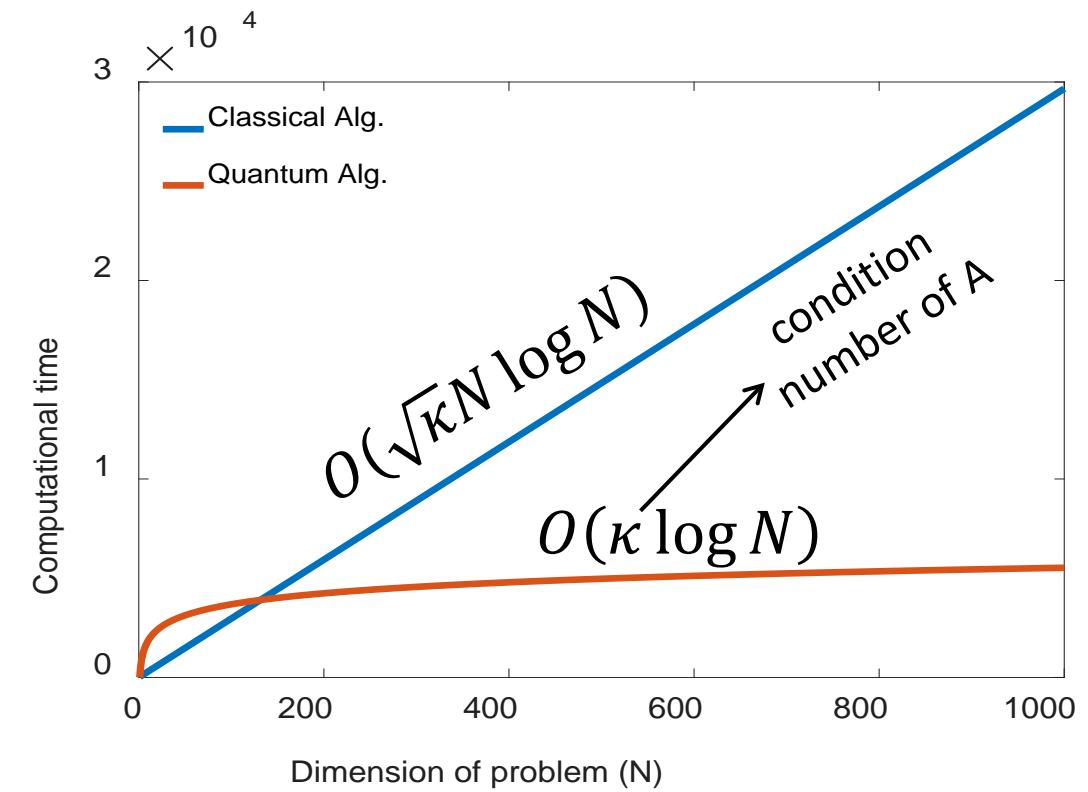
Harrow, et al. PRL, 2009

$$A \vec{x} = b$$

↓

$$|\vec{x}\rangle = A^{-1} |b\rangle_{N\text{-qubit}}$$

A N × N Matrix
 \vec{x} Unknown vector
 b Unit vector



Applications

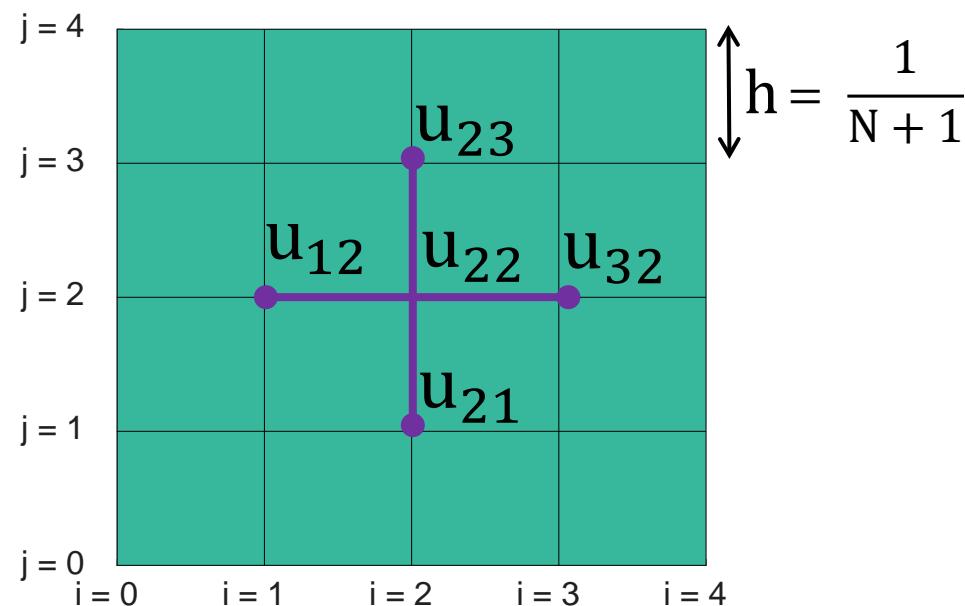
Finite Difference Modeling

Modeling in the frequency domain

$$\nabla^2 u(\vec{r}, \omega) + \frac{\omega^2}{c^2} u(\vec{r}, \omega) = f(\vec{r}, \omega)$$

Discretization

$$u_{ij} = u(ih, jh)$$
$$f_{ij} = f(ih, jh)$$



$$A \begin{bmatrix} & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \\ \vdots \\ \vdots \\ u_{N \times N} \end{bmatrix} = h^2 \begin{bmatrix} f_1 \\ f_2 \\ \vdots \\ \vdots \\ f_{N \times N} \end{bmatrix} = b$$

Applications

Finite Difference Modeling

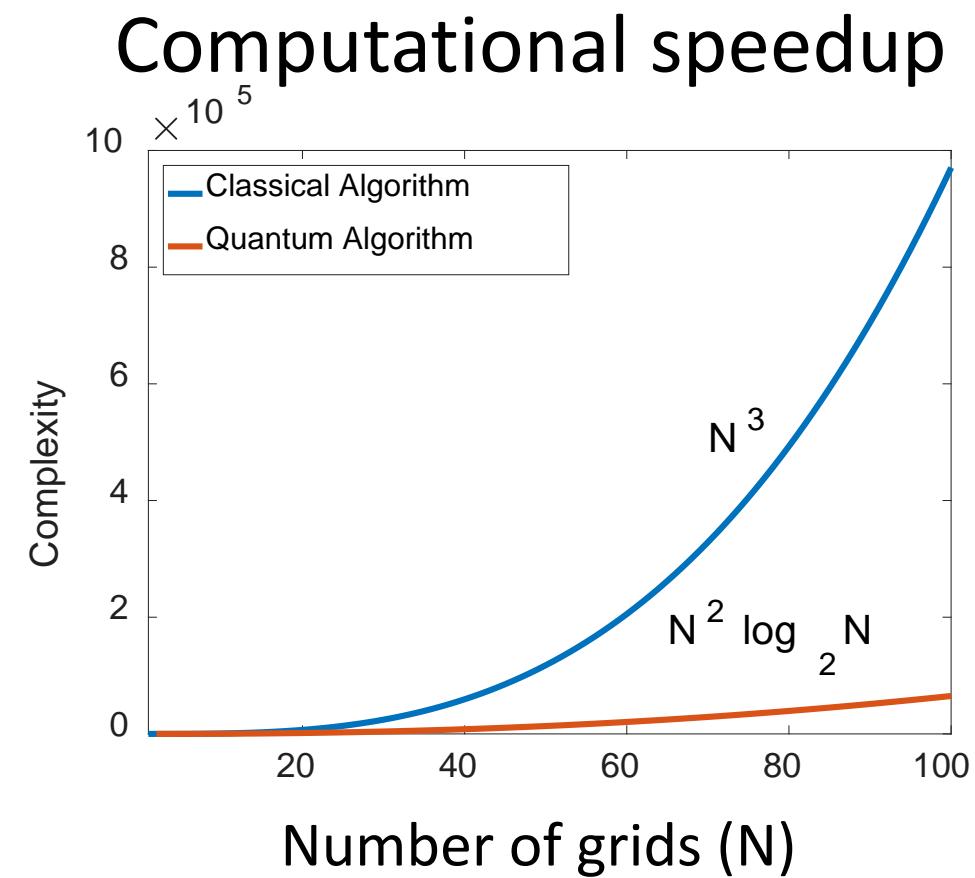
Classical algorithm
running time

$$O(n_s n_t N^3)$$

N Number of grids
 n_s Number of shots
 n_t Number of time steps

Quantum Algorithm
running time

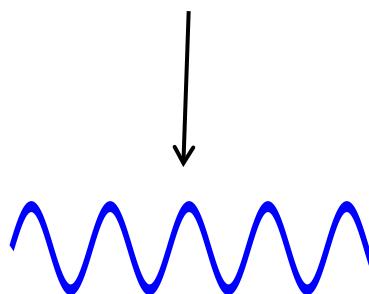
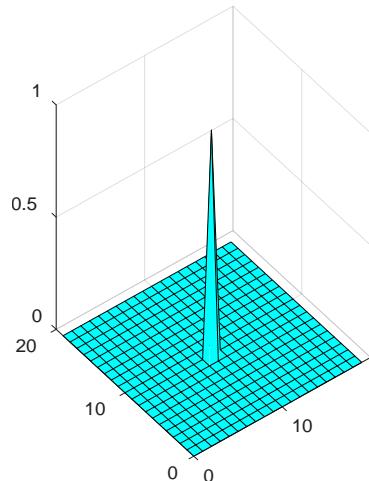
$$O(\text{poly}[\log(n_s), \log(n_t)] N^2 \log N)$$



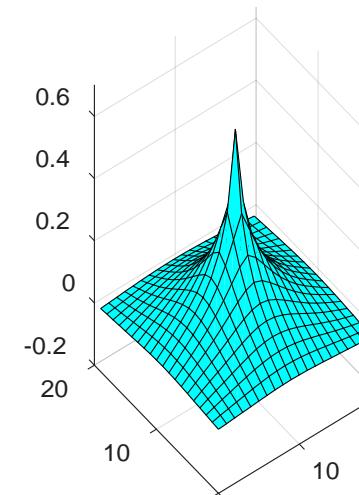
Applications

Quantum Simulator for Modeling

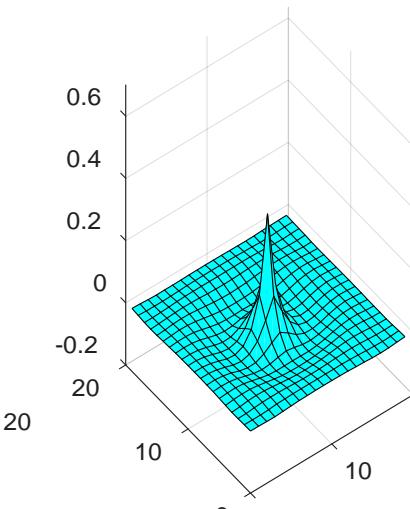
Source: right hand side



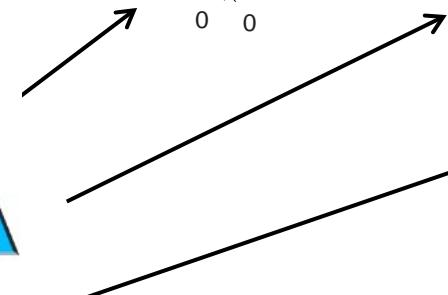
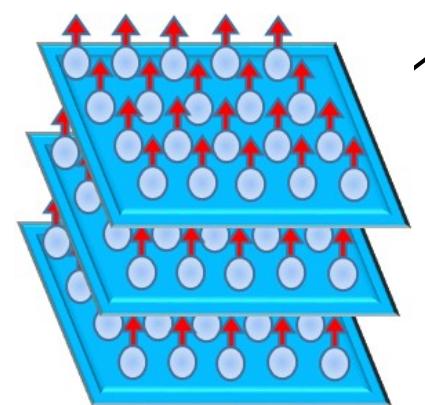
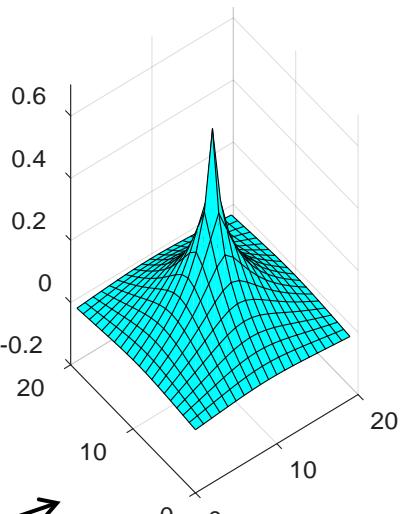
PCG method



LSQR method



Exact solution



Summary

- Quantum computing offers the exponential speedup in terms of number of grids, time steps and number of shots in seismic wave modeling
- Develop a quantum computing software package for seismic applications

Future work

- Quantum algorithms for seismic migration and inversion (imaging)
- Design and propose a quantum simulator for modeling and inversion

Acknowledgement

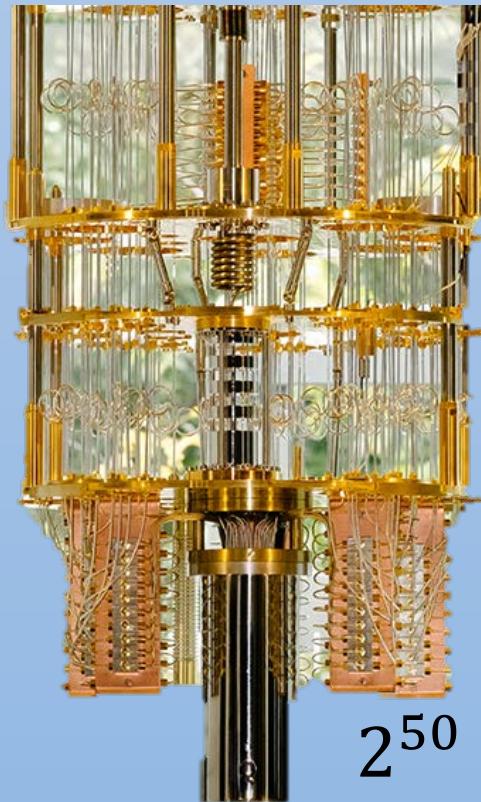
- Dr. Kris Innanen
- Dr. Hassan Khaniani
- Dr. Sam Gray
- CREWES sponsors and staff
- NSERC

Thank you

Quantum computer

IBMQ (2017)

50 qubit quantum computer



$$2^{50} \approx 10^{17} \text{ Flops}$$

Blue Gene (IBM)

most powerful supercomputer



$$2^{45} \approx 10^{15} \text{ Flops}$$

Quantum computer

Quantum computation must be done on a time-scale
less than the **decoherence time**

What is important → (decoherence time/gate operation time)

decoherence time $\sim 10^{-6}$ s

Typical gate operation time $\sim 10^{-12}$ s



$$\frac{10^{-6}\text{s}}{10^{-12}\text{s}} = 10^6$$

Number of operation that
can be executed before
the Quantum state decays

To factor $21 = 7 \times 3$ we need 10^5 gates